

CS 360 — Introduction to the Theory of Computing

Assignment 5

University of Waterloo, Spring 2018

Due 5:00 PM, July 25, 2018.

1. A *useless state* is a state that is never entered on any input string. Show that the language

$$L = \{\langle M \rangle \mid M \text{ is a Turing machine and has a useless state}\}$$

is undecidable.

2. Show that the reduction relation \leq is transitive.
3. Show that if $\mathbf{P} = \mathbf{NP}$, then every language $A \in \mathbf{P} \setminus \{\emptyset, \Sigma^*\}$ is \mathbf{NP} -complete.
4. Recall that the edit distance $d(x, y)$ of two words x and y is the fewest number of insertions, deletions, and substitutions required to transform x into y . We can extend this definition to the edit distance between two languages L_1 and L_2 by

$$d(L_1, L_2) = \min\{d(x, y) \mid x \in L_1, y \in L_2\}.$$

In other words, the distance between two languages is the distance between their two closest words.

Show that if L_1 and L_2 are context-free languages, then $d(L_1, L_2)$ is uncomputable.

5. Let the class \mathbf{coNP} be defined by

$$\mathbf{coNP} = \{L \mid \bar{L} \in \mathbf{NP}\}.$$

That is, \mathbf{coNP} is the class of languages that are complements of languages in \mathbf{NP} . As with many problems in computational complexity, it is not known whether $\mathbf{NP} = \mathbf{coNP}$.

Show that $\mathbf{NP} = \mathbf{coNP}$ if and only if there exists a language L such that $L \in \mathbf{coNP}$ and L is \mathbf{NP} -complete.